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NewSQL Overview

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About myself

- ❖ MIPT
- ❖ MCST, Elbrus compiler project
- ❖ Echo, real-time social platform (PaaS)
- ❖ DevZen podcast (<http://devzen.ru>)



History of SQL

- ❖ Relational Model in 1970

- ❖ disk-oriented

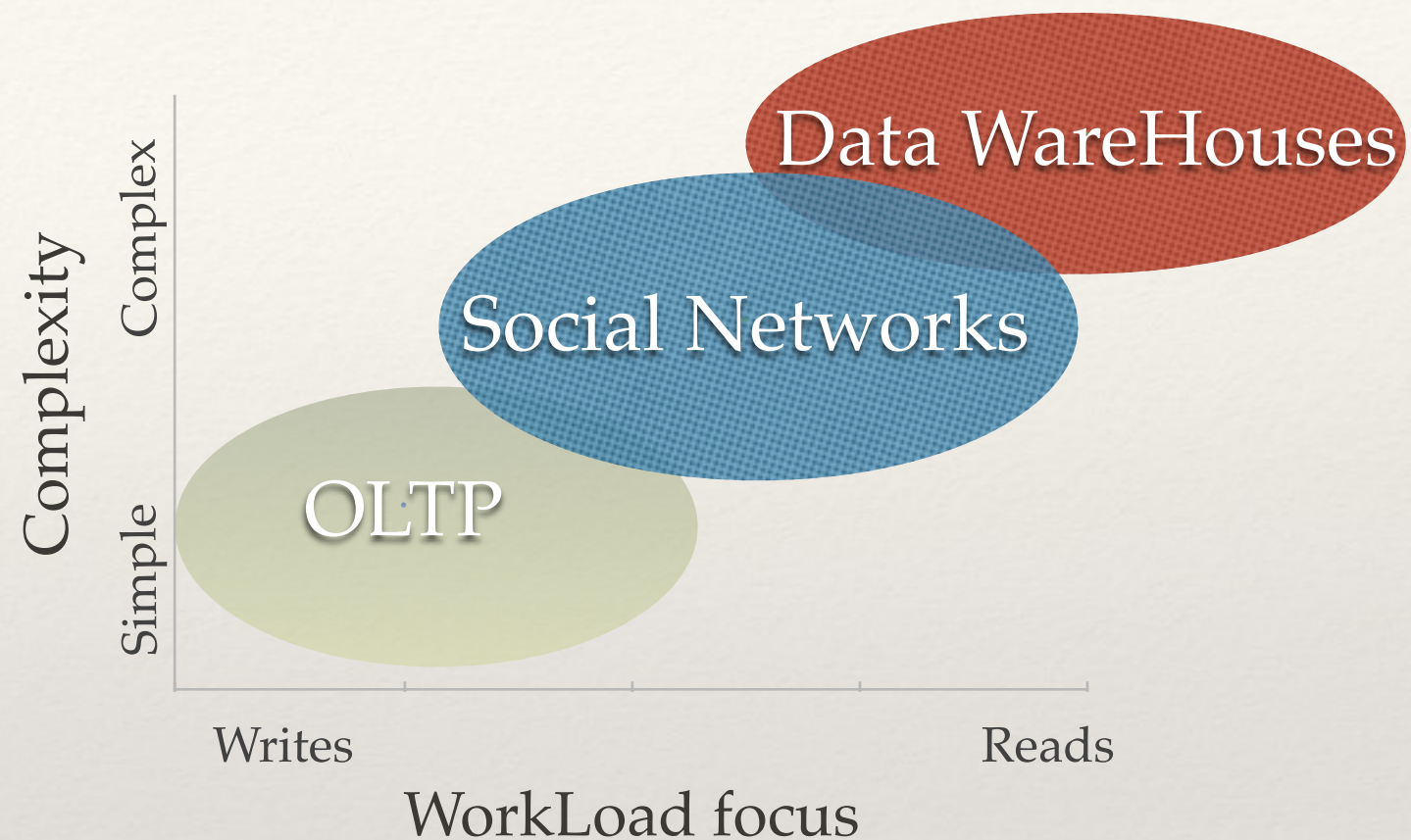
- ❖ rows

- ❖ sql

- ❖ “One size fits all” doesn’t work:

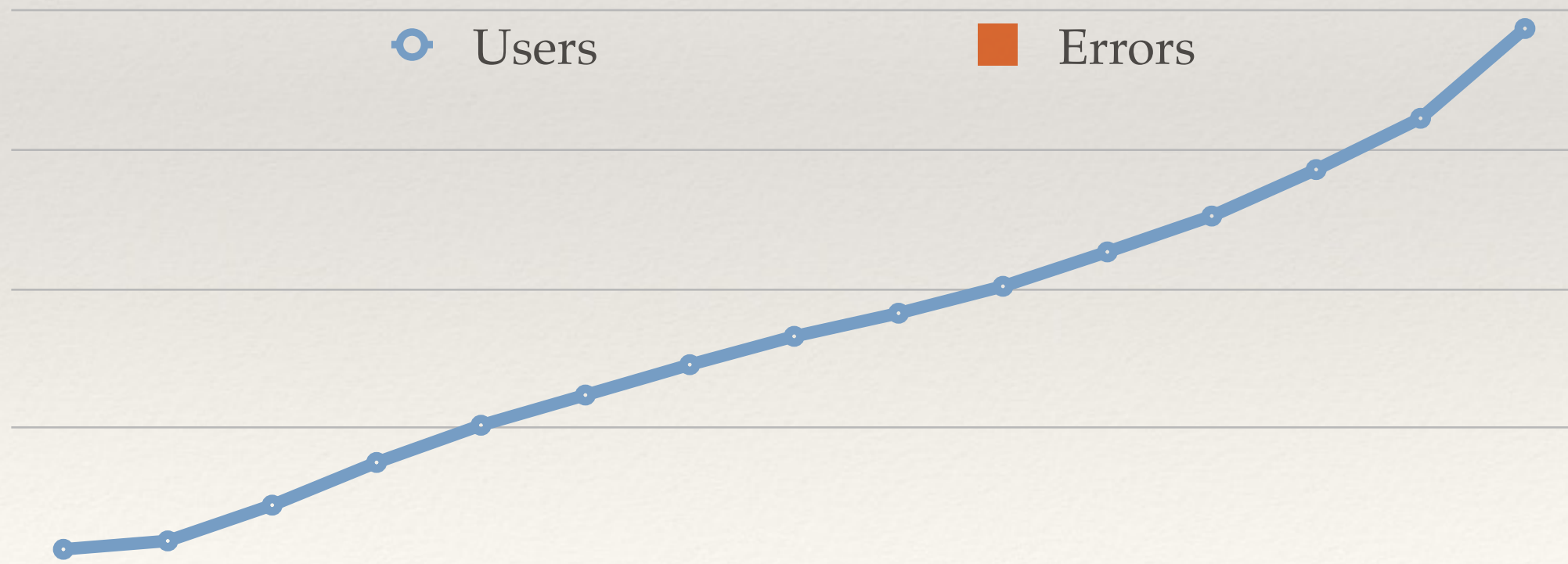
- ❖ Column-oriented data warehouses for OLAP.

- ❖ Key-Value storages, Document storages



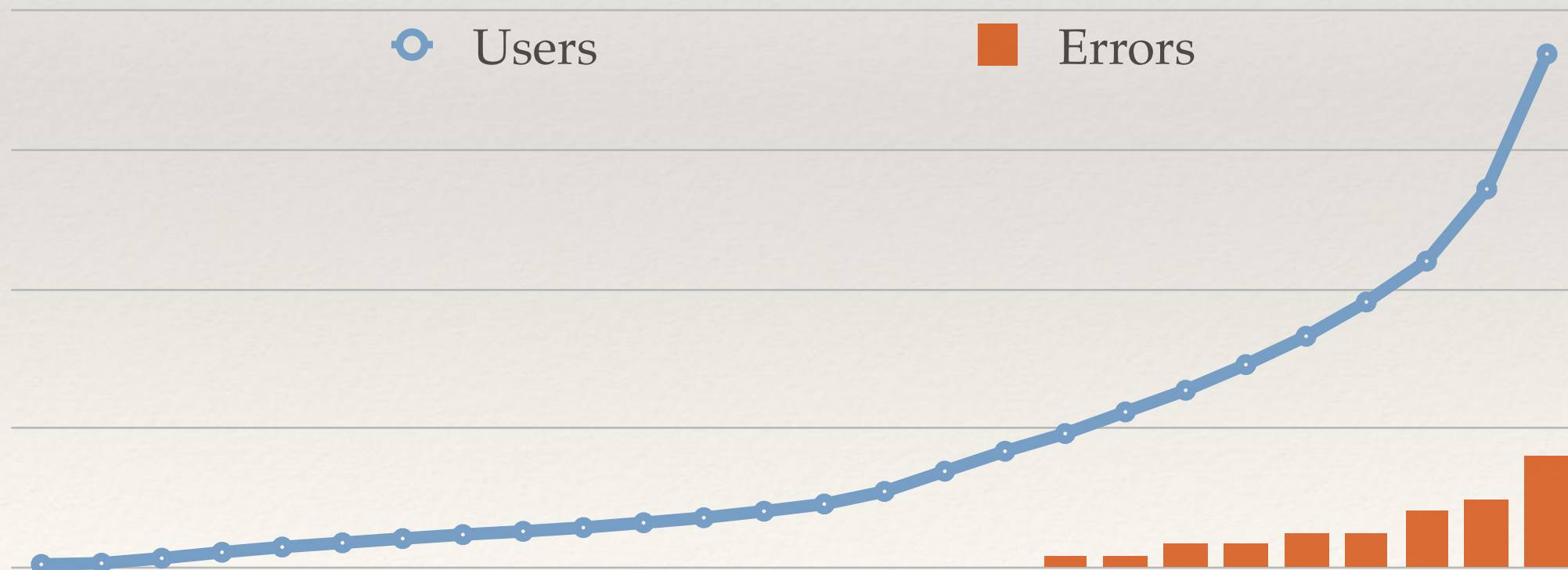
Startups lifecycle

- ❖ Start: no money, no users, open source



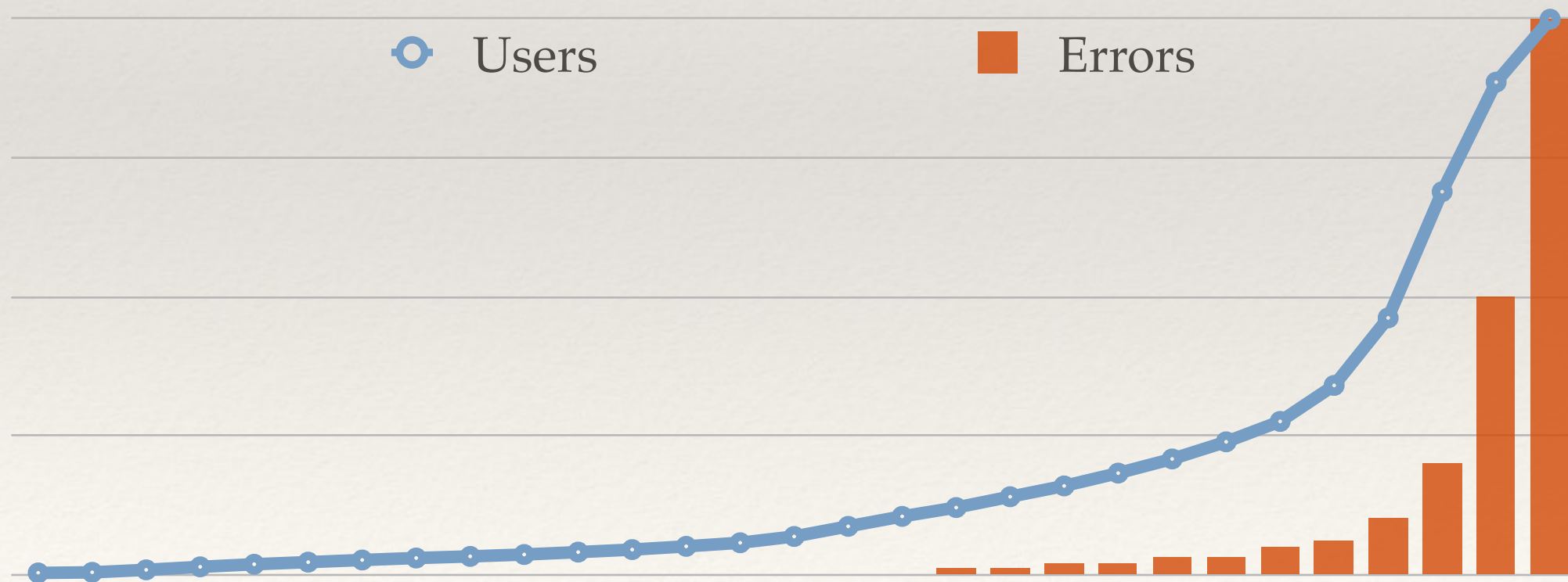
Startups lifecycle

- ❖ Start: no money, no users, open source
- ❖ Middle: more users, storage optimization



Startups lifecycle

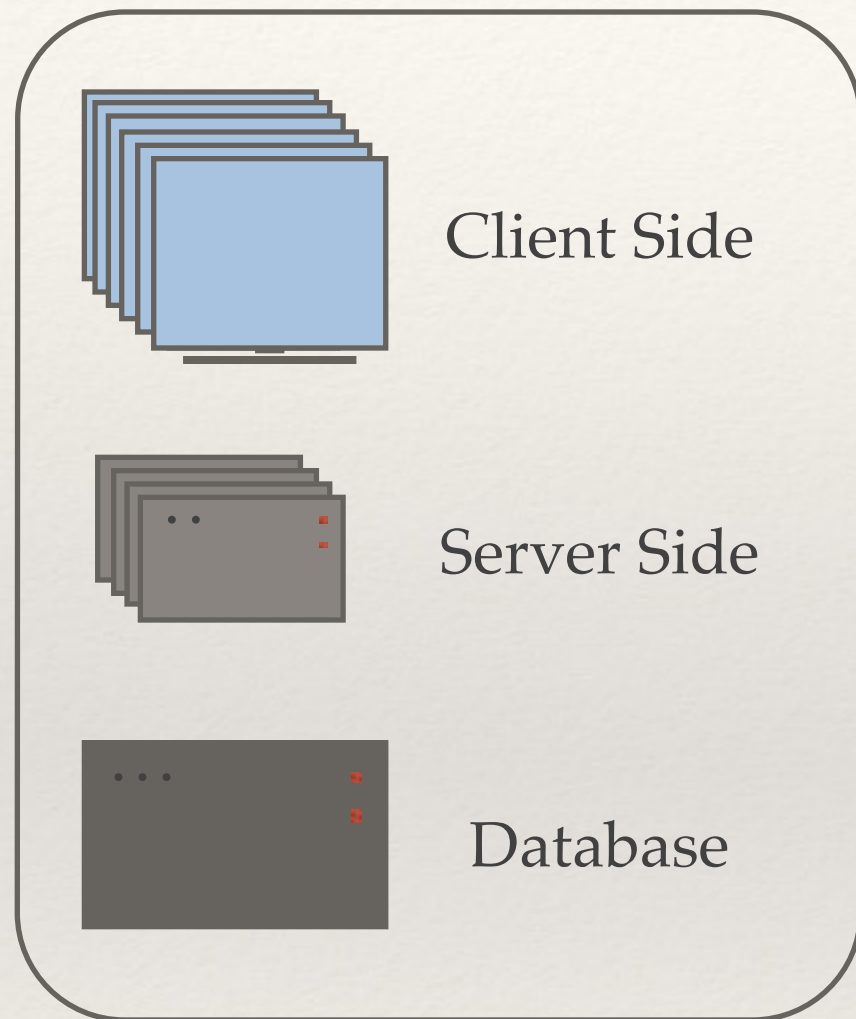
- ❖ Start: no money, no users, open source
- ❖ Middle: more users, storage optimization
- ❖ Final: plenty of users, storage failure



New requirements

- ❖ Large scale systems, with huge and growing data sets
 - ❖ 9M messages per hour in Facebook
 - ❖ 50M messages per day in Twitter
- ❖ Information is frequently generated by devices
- ❖ High concurrency requirements
- ❖ Usually, data model with some relations
- ❖ Often, transactional integrity

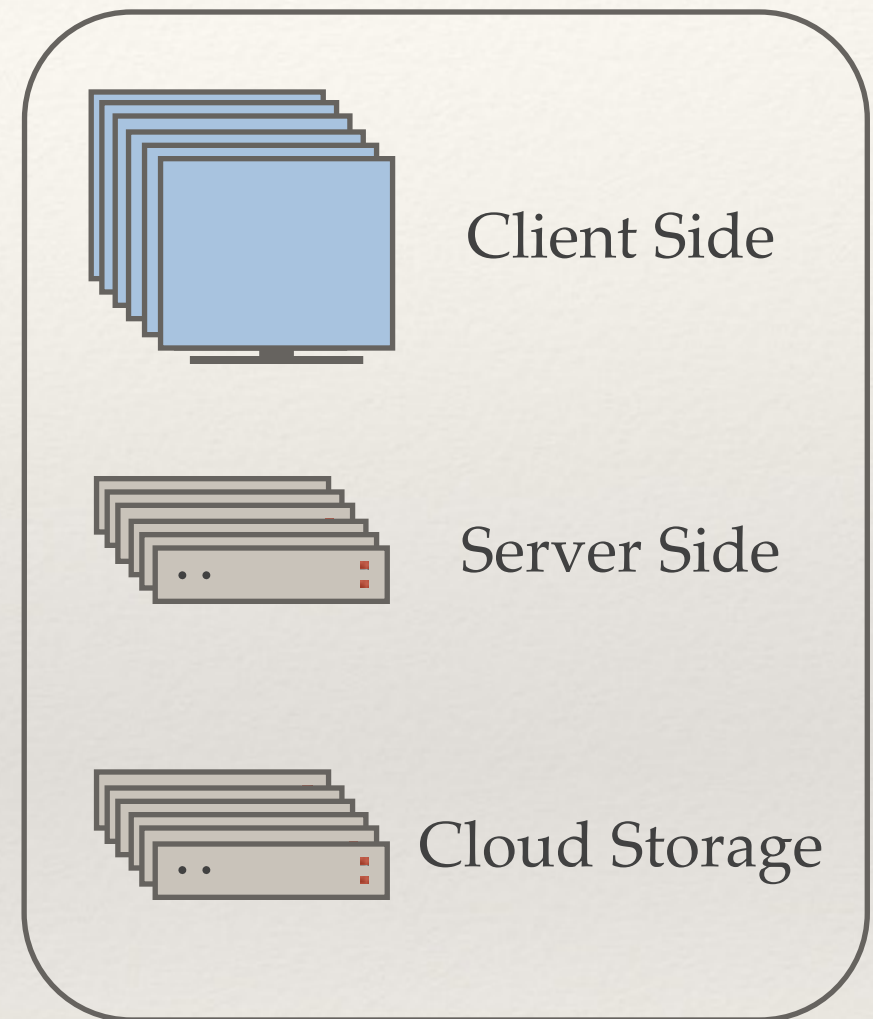
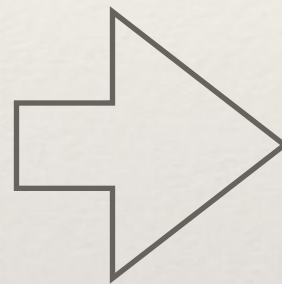
Trends: architecture change



Consistency, transactions: Database

Storage optimization: Database

Scalability: Client Side



Consistency, transactions: Cloud

Storage optimization: Cloud

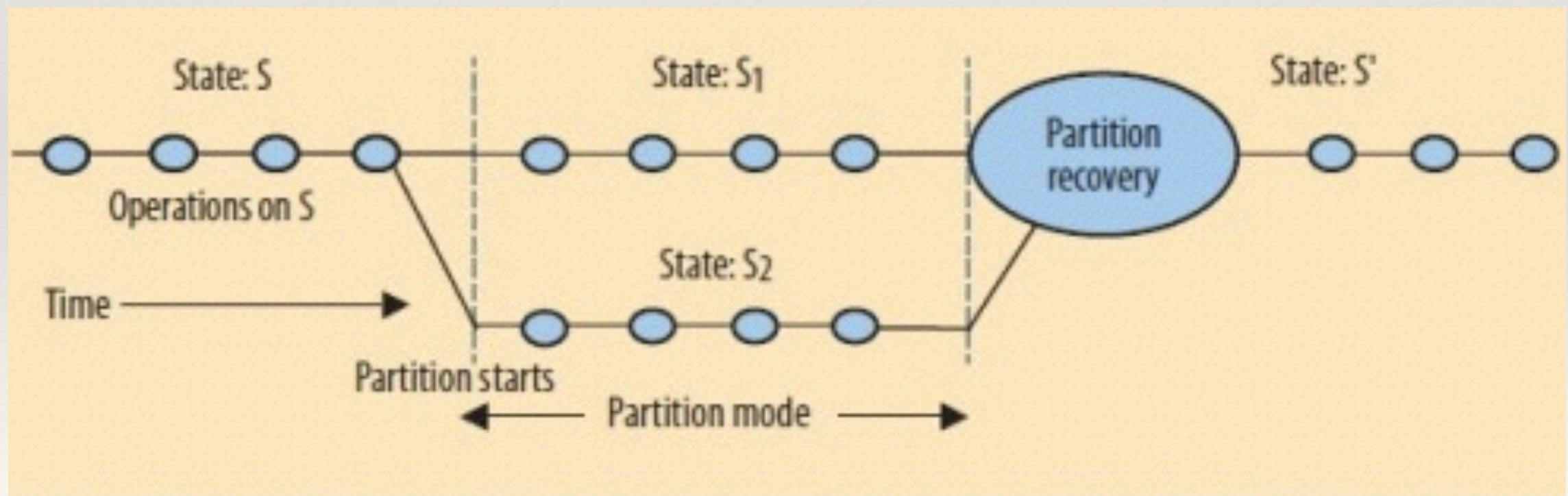
Scalability: All levels

Trends: architecture change

- ❖ CAP: **c**onsistency, **a**vailability, **p**artitioning
- ❖ ACID: **a**tomicity, **c**onsistency, **i**solation, **d**urability
- ❖ BASE: **b**asically **a**vailable, **s**oft state, **e**ventual consistency

Trends: architecture change

- ❖ 'P' in CAP is not discrete
- ❖ Managing partitions: detection, limitations in operations, recovery



NoSQL

- ❖ CAP: first 'A', then 'C': finer control over availability
- ❖ Horizontal scaling
- ❖ Not a “relational model”, custom API
- ❖ Schemaless
- ❖ Types: Key-Value, Document, Graph, ...

Application-level sharding

- ❖ Additional application-level logic
- ❖ Difficulties with cross-sharding transactions
- ❖ More servers to maintain
- ❖ More components — higher prob for breakdown

NewSQL: definition

“A DBMS that delivers the scalability and flexibility promised by NoSQL while retaining the support for SQL queries and/or ACID, or to improve performance for appropriate workloads.”

451 Group

NewSQL: definition

- ❖ *SQL as the primary interface*
- ❖ *ACID support for transactions*
- ❖ *Non-locking concurrency control*
- ❖ *High per-node performance*
- ❖ *Scalable, shared nothing architecture*

Michael Stonebraker

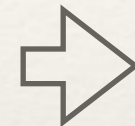
Shared nothing architecture

- ❖ No single point of failure
- ❖ Each node is independent and self-sufficient
- ❖ No shared memory or disk
- ❖ Scale infinitely
- ❖ Data partitioning
- ❖ Slow multi-shards requests

Column-oriented DBMS

- ❖ Store content by column rather than by row

John	Smith	20
Joe	Smith	30
Alice	Adams	50



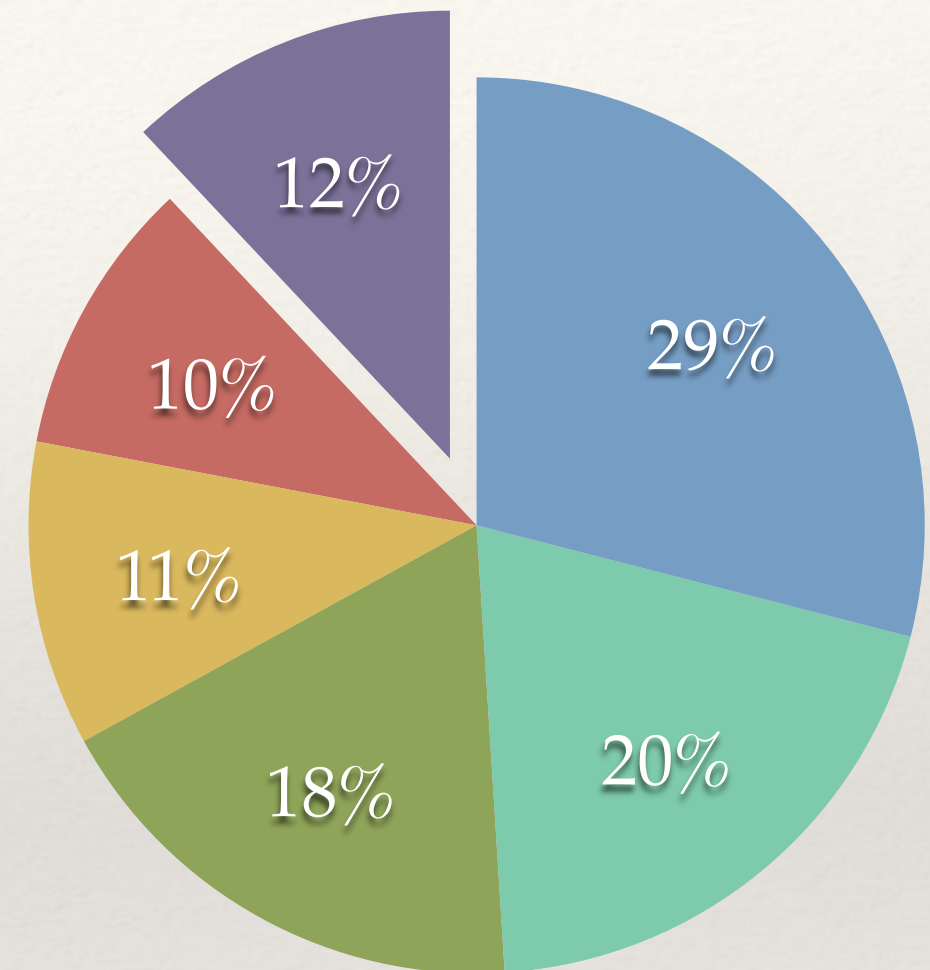
John:001; Joe:002; Alice:003.
Smith:001,002; Adams:003.
20:001; 30:002; 50:003.

- ❖ Efficient in hard disk access
- ❖ Good for sparse and repeated data
- ❖ Higher data compression
- ❖ More reads / writes for large records with a lot of fields
- ❖ Better for relatively infrequent writes, lots of data throughput on reads (OLAP, analytic requests).

Traditional DBMS overheads

by Stonebraker & research group

- Buffer Management
- Logging
- Locking
- Index management
- Latching
- Useful work



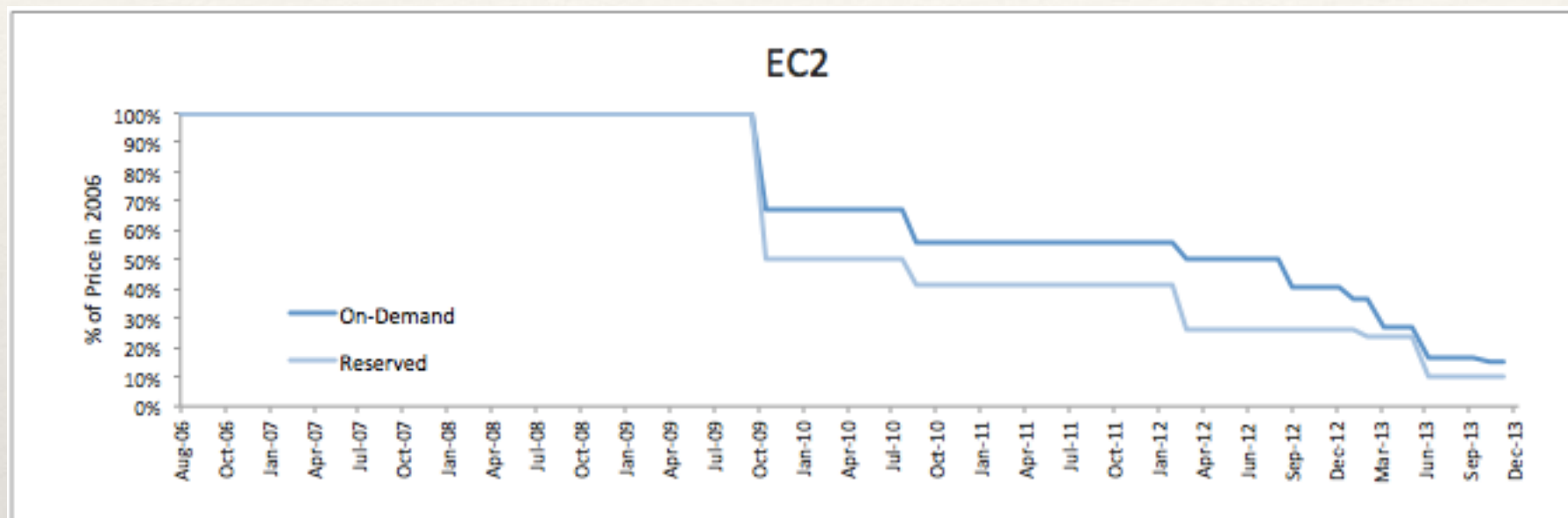
“Removing those overheads and running the database in main memory would yield orders of magnitude improvements in database performance”

In-memory storage

- ❖ High throughput
- ❖ Low latency
- ❖ No Buffer Management
- ❖ If serialized, no Locking or Latching

In-memory storage: price

Amazon price reduction



Current price for 1TB (~4 instances of 'r3.8xlarge' type)

	on-demand	3Y-reserved plan
per hour	11.2 \$	3.9 \$
per month	8.1K \$	2.8K \$
per year	97K \$	33,7K \$

NewSQL: categories

- ❖ New approaches: VoltDB, Clustrix, NuoDB
- ❖ New storage engines: TokuDB, ScaleDB
- ❖ Transparent clustering: ScaleBase, dbShards

NuoDB

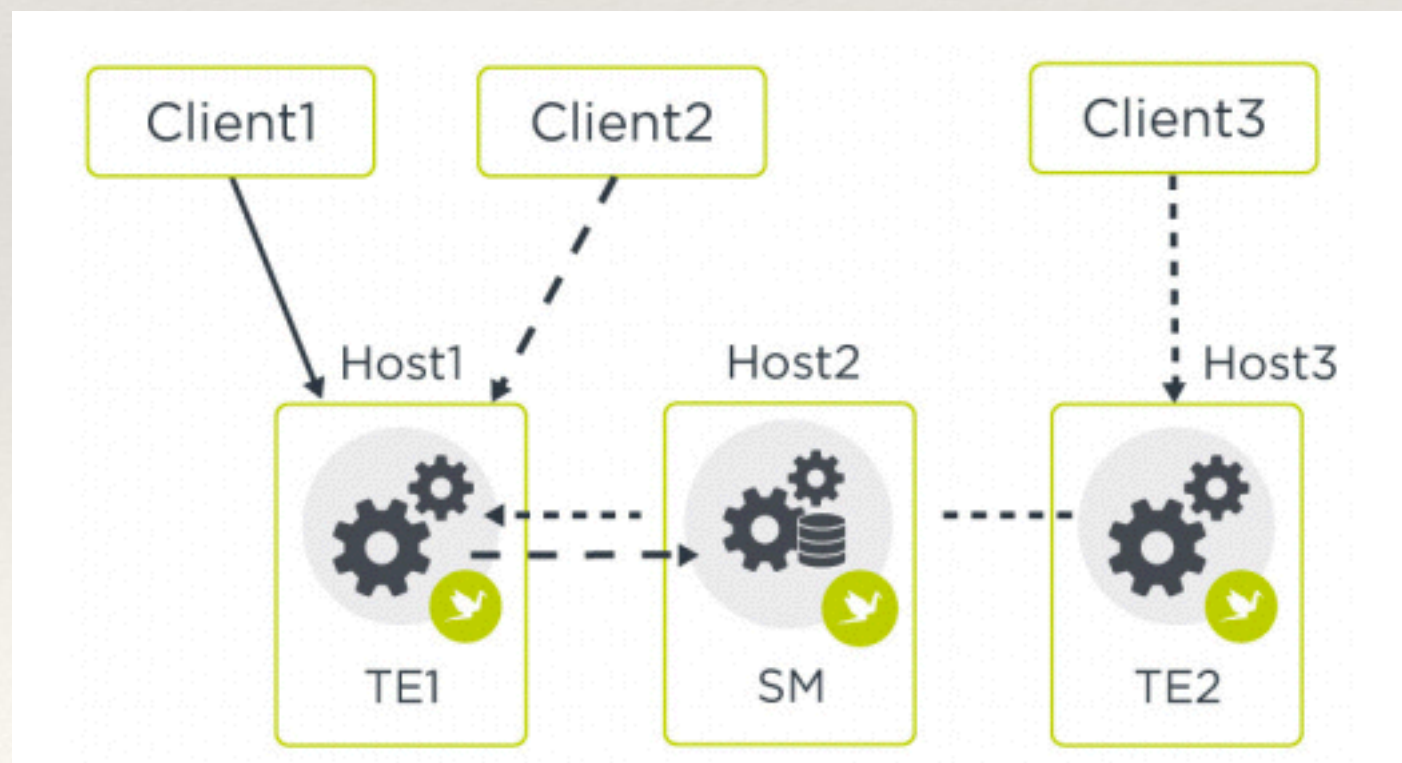


- ❖ Multi-tier architecture:
 - ❖ Administrative: managing, stats, cli, web-ui
 - ❖ Transactional: ACID except 'D', cache
 - ❖ Storage: key-value store ('D' from ACID)

NuoDB



- ❖ Everything is an 'Atom'
- ❖ Peer-to-peer communication, encrypted sessions
- ❖ MVCC + Append-only storage

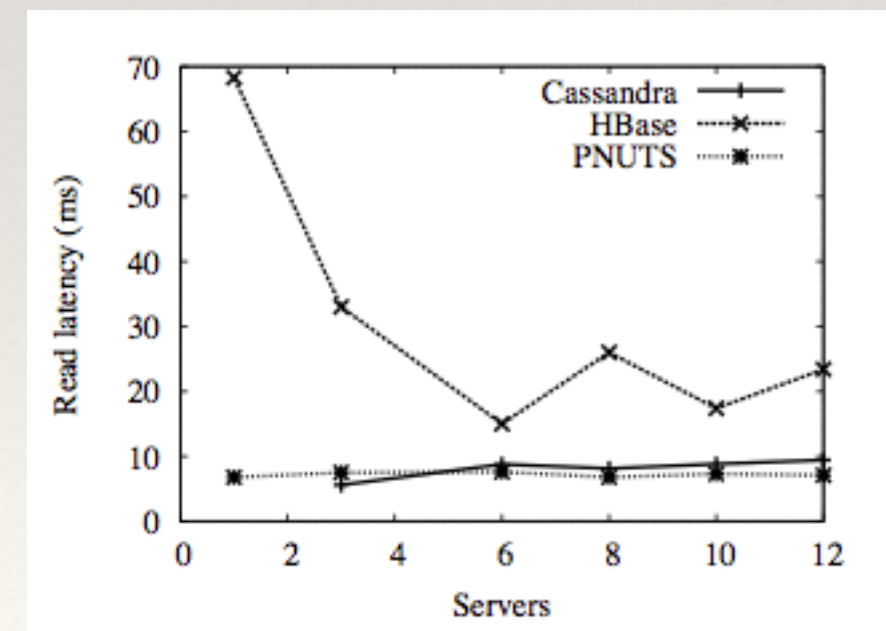
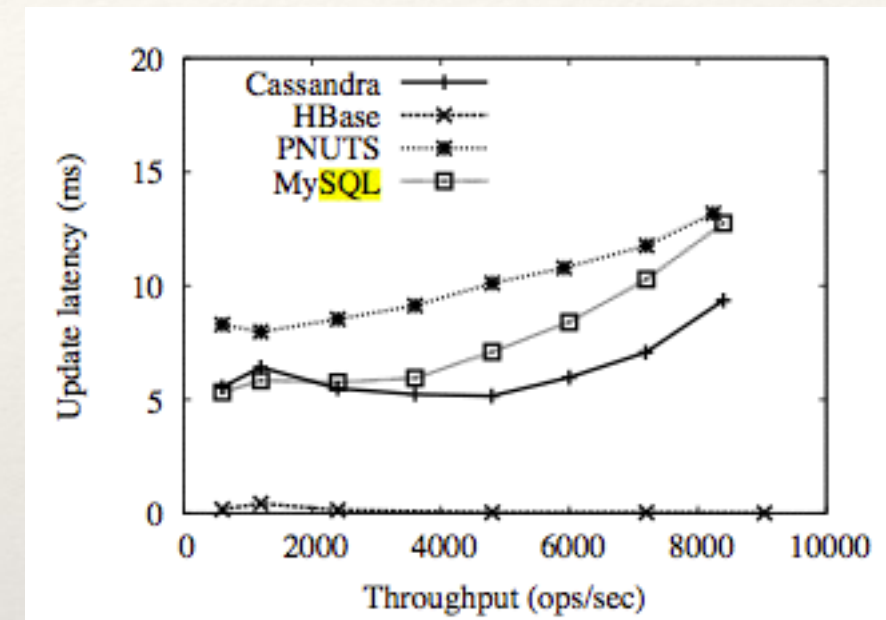


NuoDB: CAP & ACID

- ❖ `CP` system. Need majority of nodes to work
- ❖ If split to two equal parts -> stop
- ❖ Several consistency modes including 'consistent_read'

YCSB

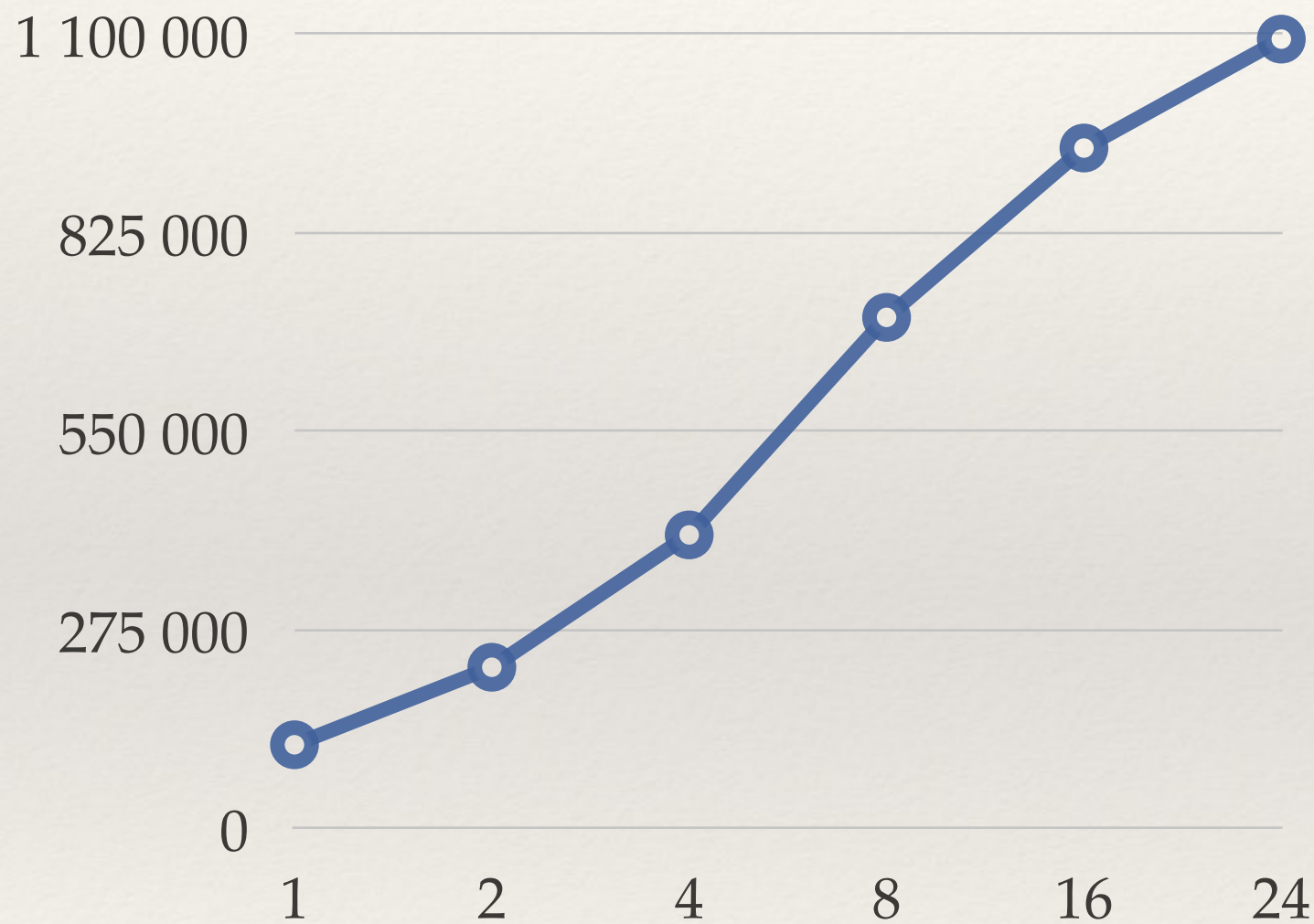
- ❖ Yahoo Cloud Serving Benchmark
- ❖ Key-value: insert / read / update / scan
- ❖ Measures:
 - ❖ Performance: latency / throughput
 - ❖ Scaling: elastic speedup



NuoDB: YCSB



Throughput, tps / nodes



Update latency, μ s



Read latency, μ s



5% updates, 95% reads

Hosts: 32GB, Xeon 8 cores, 1TB HDD, 1Gb LAN

VoltDB

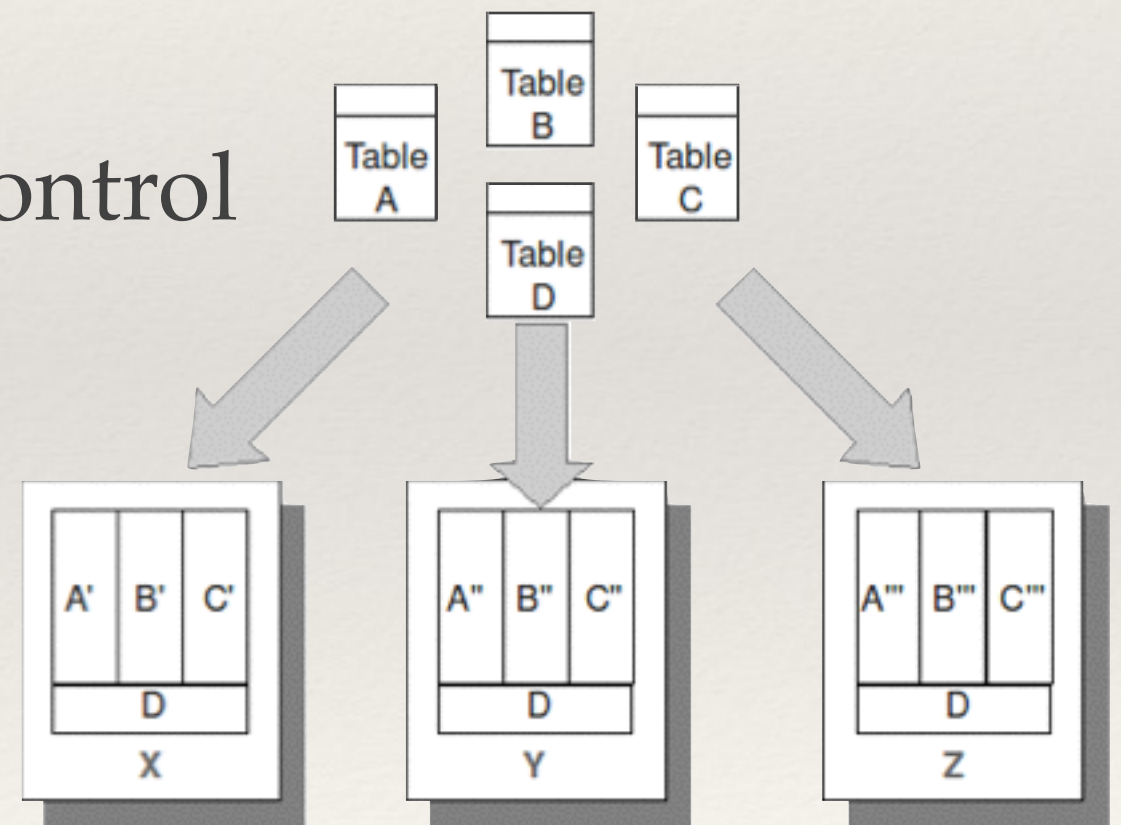


- ❖ In-memory storage
- ❖ Stored procedure interface, async / sync proc execution
- ❖ Serializing all data access
- ❖ Horizontal partitioning
- ❖ Multi-master replication (“K-safety”)
- ❖ Snapshots + Command Logging

VoltDB



- ❖ Open-source, community edition is under GPLv3.
- ❖ Java + C++
- ❖ Partitioning and Replication control

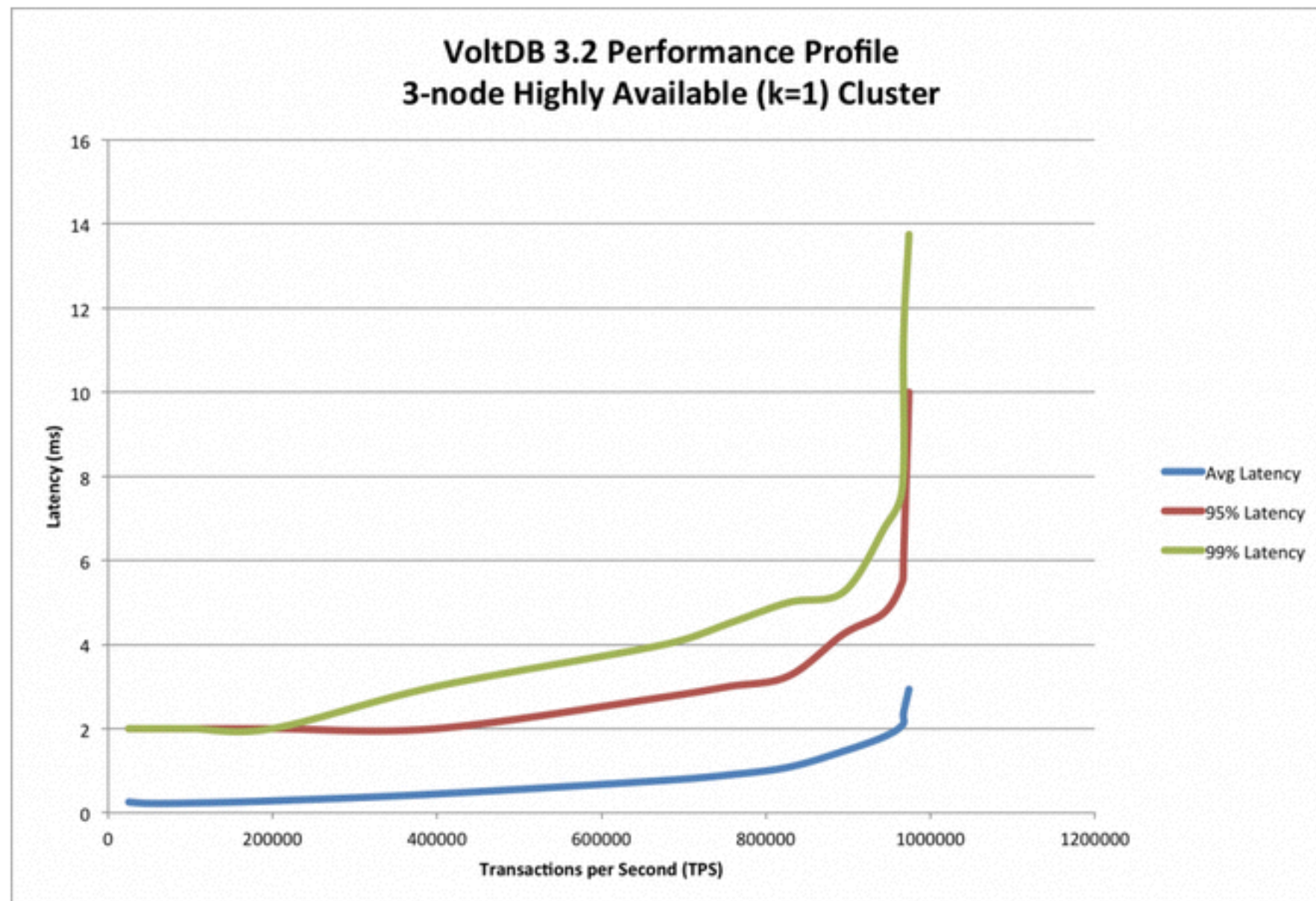


VoltDB: CAP & ACID



- ❖ Without K-safety, any node fail break the whole DB
- ❖ Snapshot and shutdown minor segments during network partitions
- ❖ Single-partition transactions are very fast
- ❖ Multi-partition transactions are slower (manager), try to avoid (1000s tps in '13, no updates since)

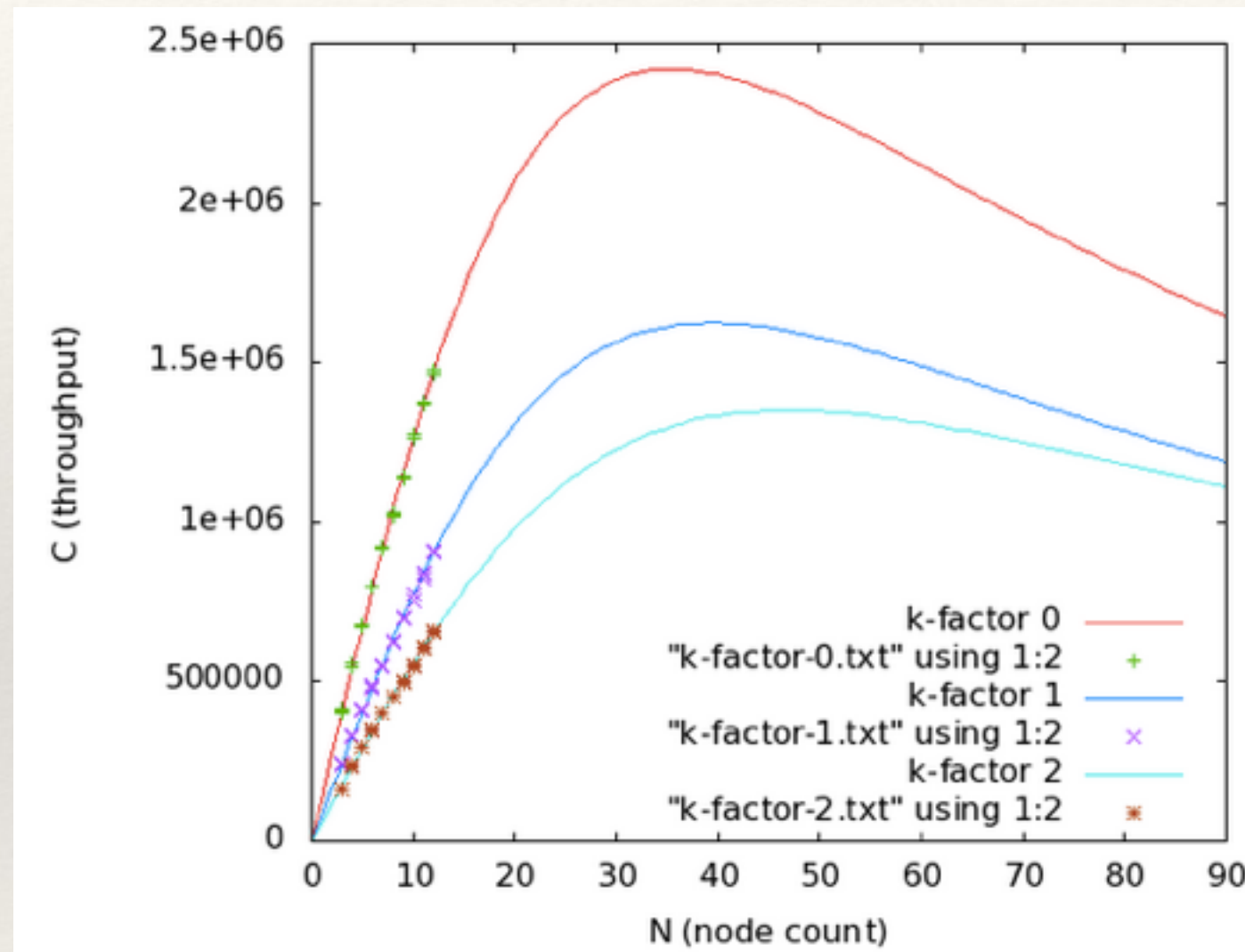
VoltDB: key-value bench



90% reads, 10% writes

3 nodes: 64GB, dual 2.93GHz intel 6 core processors

VoltDB: “voter” bench

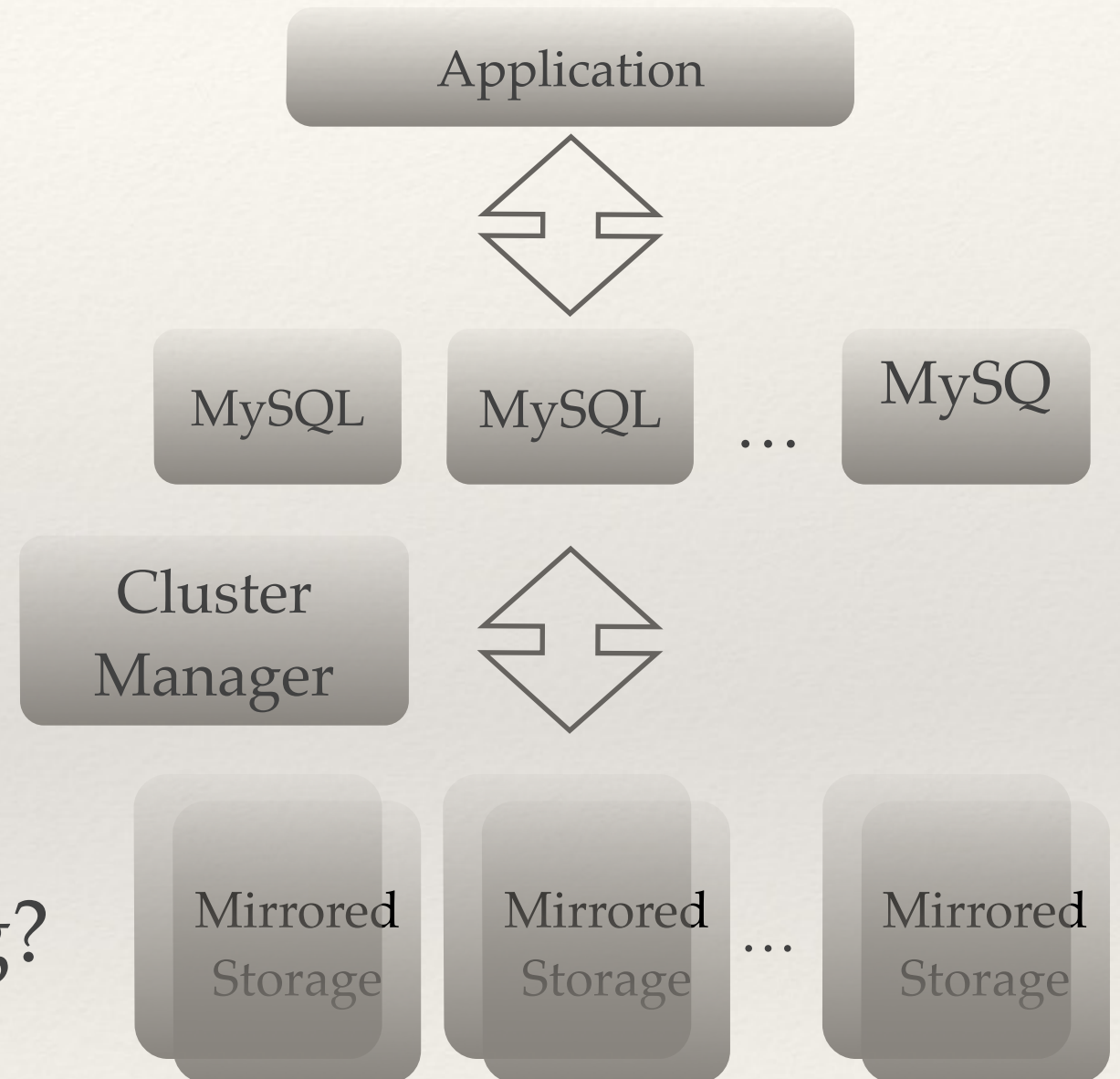


26 SQL statements per transaction

ScaleDB

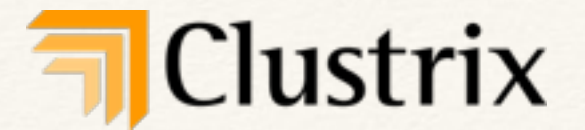


- ❖ Multi-master
- ❖ Shared data
- ❖ Cluster manager to solve conflicts (locks)
- ❖ ACID?
- ❖ Network Partition Handling?
- ❖ Scaling?

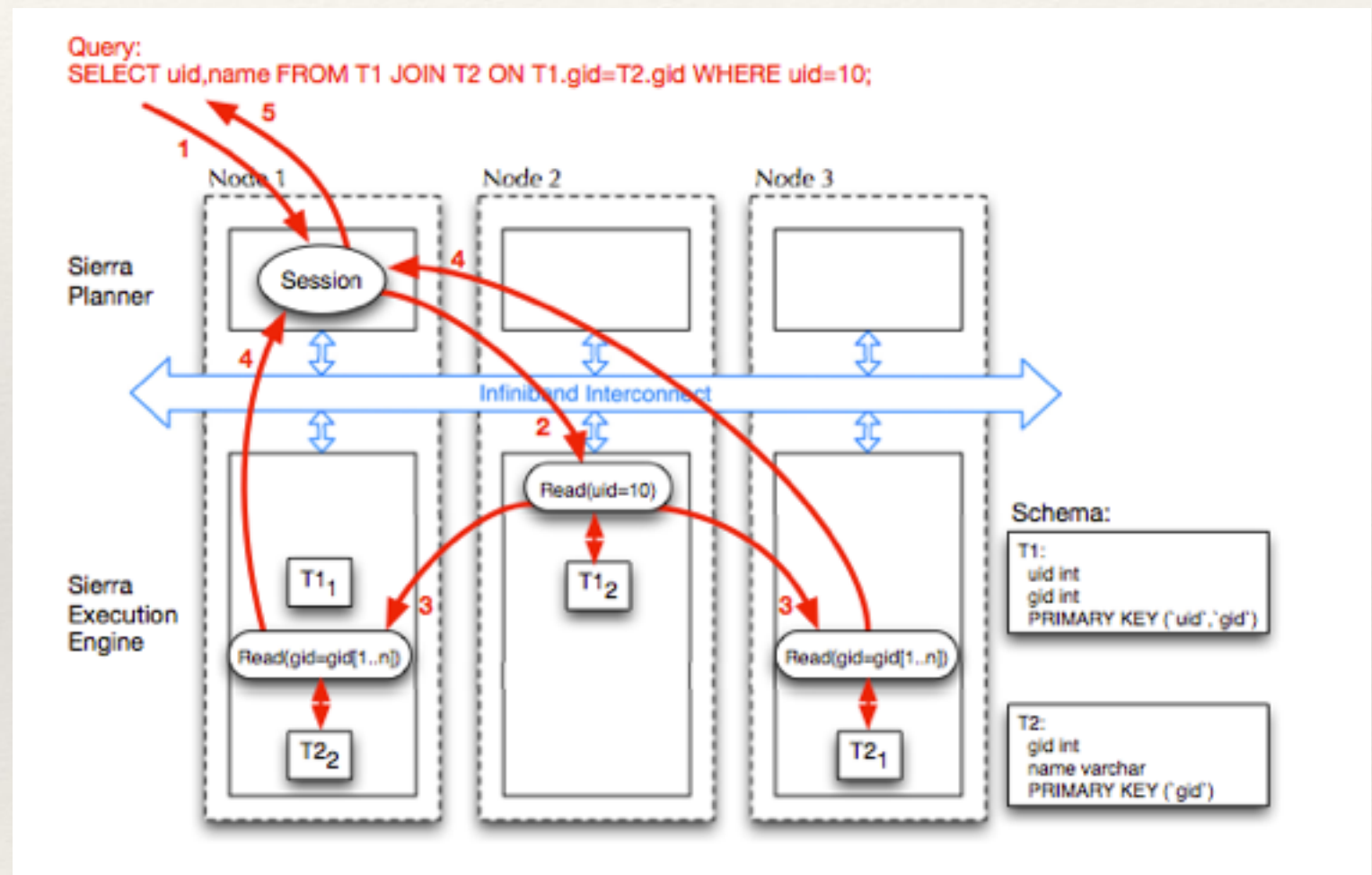


- ❖ “Query fragment” - basic primitive of the system:
 - ❖ read / write / execute function
 - ❖ modify control flow
 - ❖ perform synchronisation
 - ❖ send rows to query fragments on another nodes
- ❖ Data partitions: “slices” split and moved transparently
- ❖ Replication: master slice for reads + slave for redundancy

ClustrixDB



- ❖ “Move query to the data”
- ❖ Dynamic and transparent data layout
- ❖ Linear scale



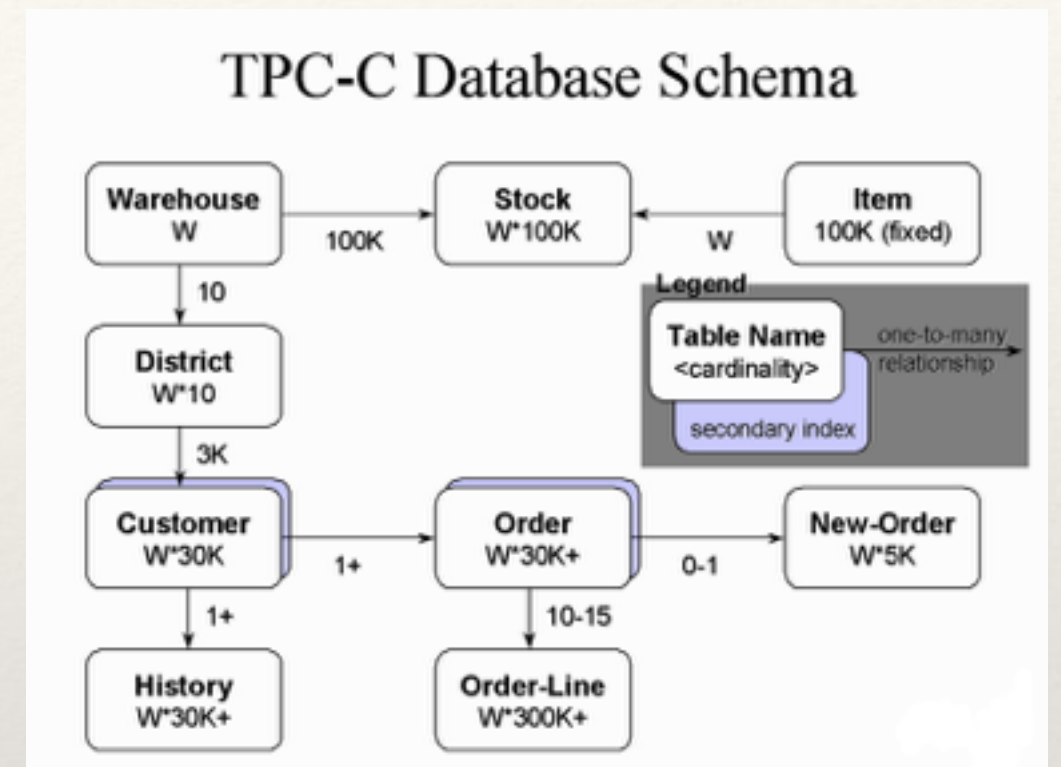
```
Schema:
T1:
  uid int
  gid int
  PRIMARY KEY ('uid','gid')
T2:
  gid int
  name varchar
  PRIMARY KEY ('gid')
```

ClustrixDB: CAP & ACID

- ❖ `CP` system. Need majority of nodes to work
- ❖ Only 'Repeatable Read' isolation level (so, 'phantom reads' are possible)
- ❖ Distributed Lock Manager for writer-writer locks (on each node)

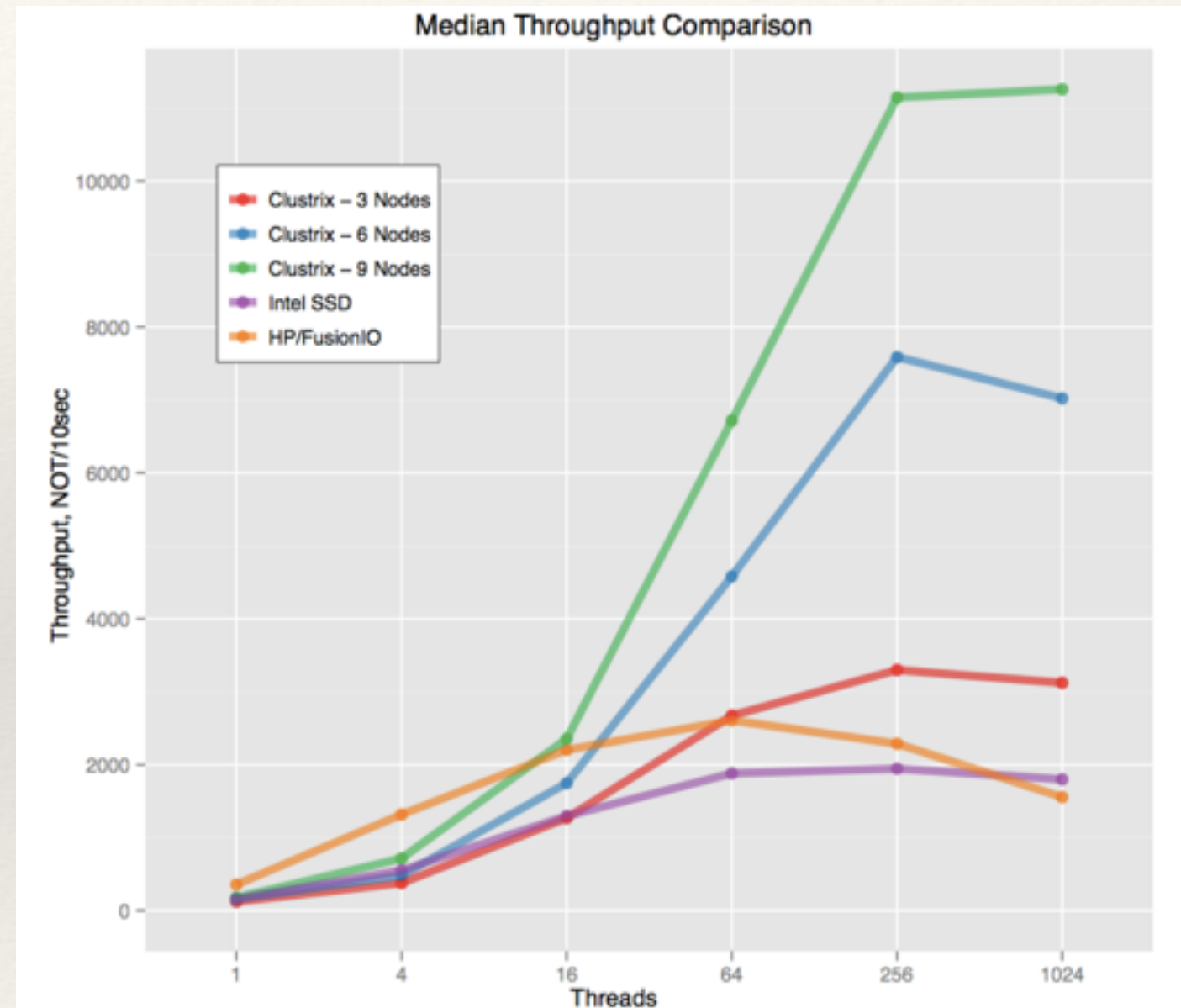
TPC-C

- ❖ Online Transaction Processing (OLTP) benchmark
- ❖ 9 types of tables
- ❖ 5 concurrent transactions of different complexity
- ❖ Productivity measured in “new-order transaction”



ClustrixDB: TPC-C Clustrix

- ❖ 5000W ~ 400GB of data
- ❖ Compared with Percona Mysql, Intel Xeon, 8 cores
- ❖ ClustrixDB nodes: “Dual 4 core Westmere processors”



ClustrixDB: example Clustrix

- ❖ 30M users, 10M logins per day
- ❖ 4.4B transactions per day
- ❖ 1.08 / 4.69 Petabytes per month writes / reads
- ❖ 42 nodes, 336 cores, 2TB memory, 46TB SSD

FoundationDB

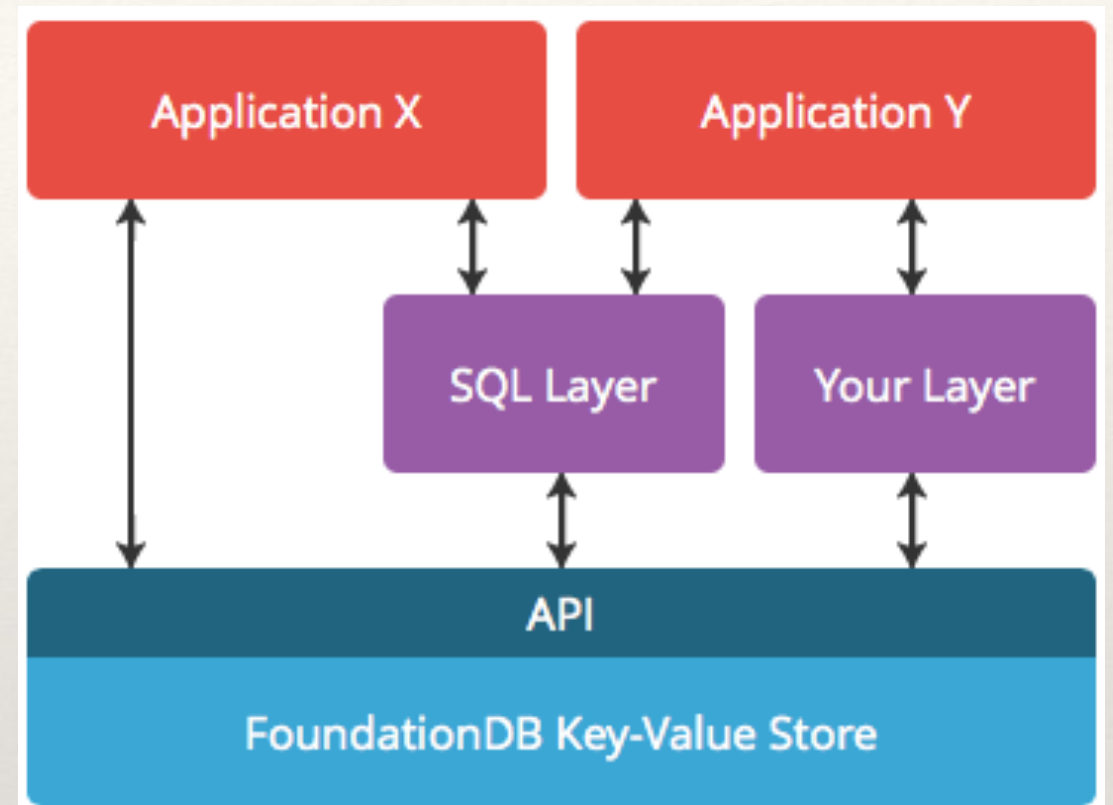


- ❖ KV store, ordered keys
- ❖ Paxos for cluster coordination
- ❖ Global ACID transactions, range operations
- ❖ Lock-free, optimistic concurrency, MVCC
- ❖ Good testing (deterministic simulation)
- ❖ Fault-tolerance (replication)
- ❖ SQL Layer (similar to Google F1 on top of Spanner)

FoundationDB



- ❖ SSD / Memory storage engine
- ❖ Layers concept
- ❖ 'CP' system with Paxos-ed coordination centres
- ❖ Written in the Flow language (translated to C++11) with actor model support
- ❖ Watches, atomic operations (e.g. 'add')



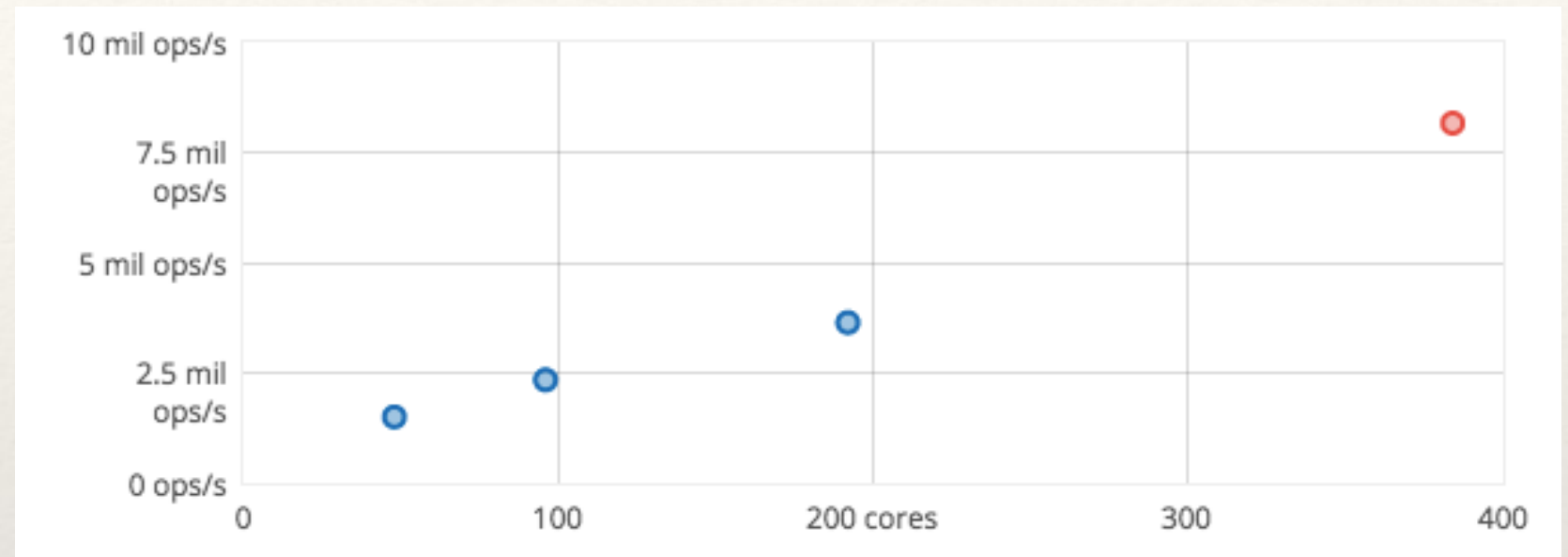
FoundationDB: CAP and ACID

- ❖ Serializable isolation with optimistic concurrency
- ❖ > 100 wps to the same key? Use another DB!
- ❖ 'CP system' (Paxos)
Need majority of coordination center to work

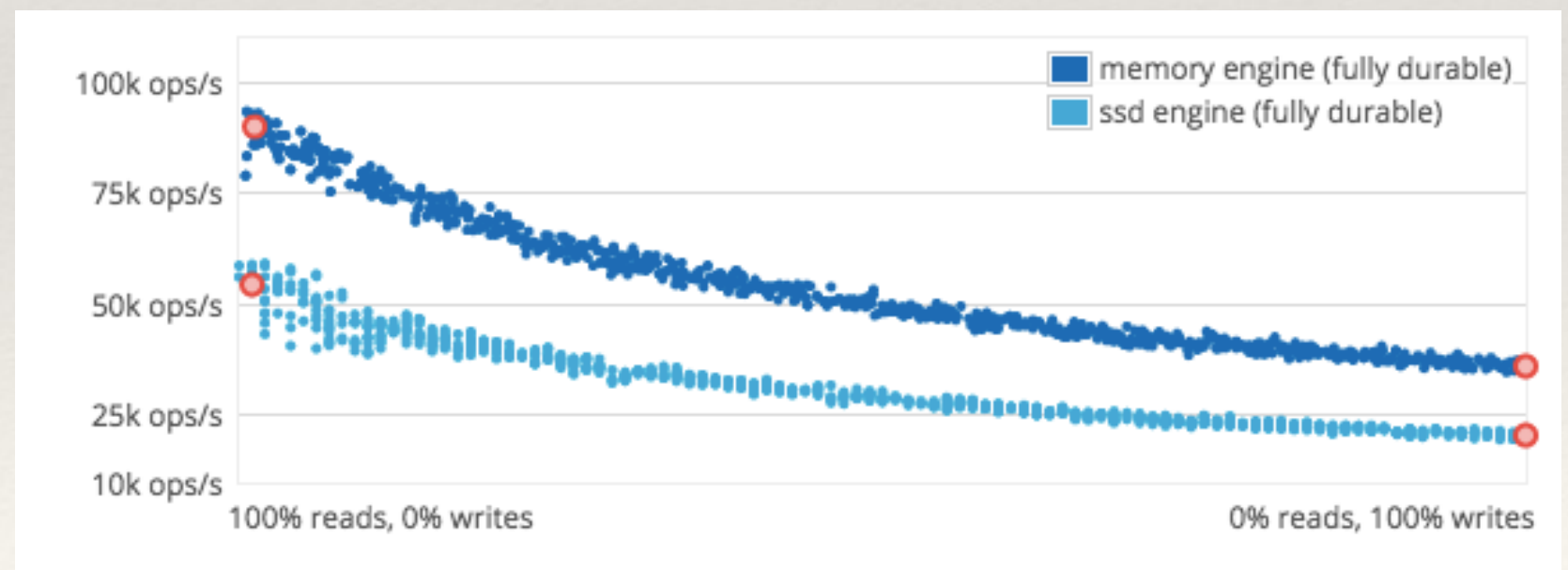
FoundationDB: KV Performance

Scaling:

up to 24 EC2 c3.8xlarge, 16 cores

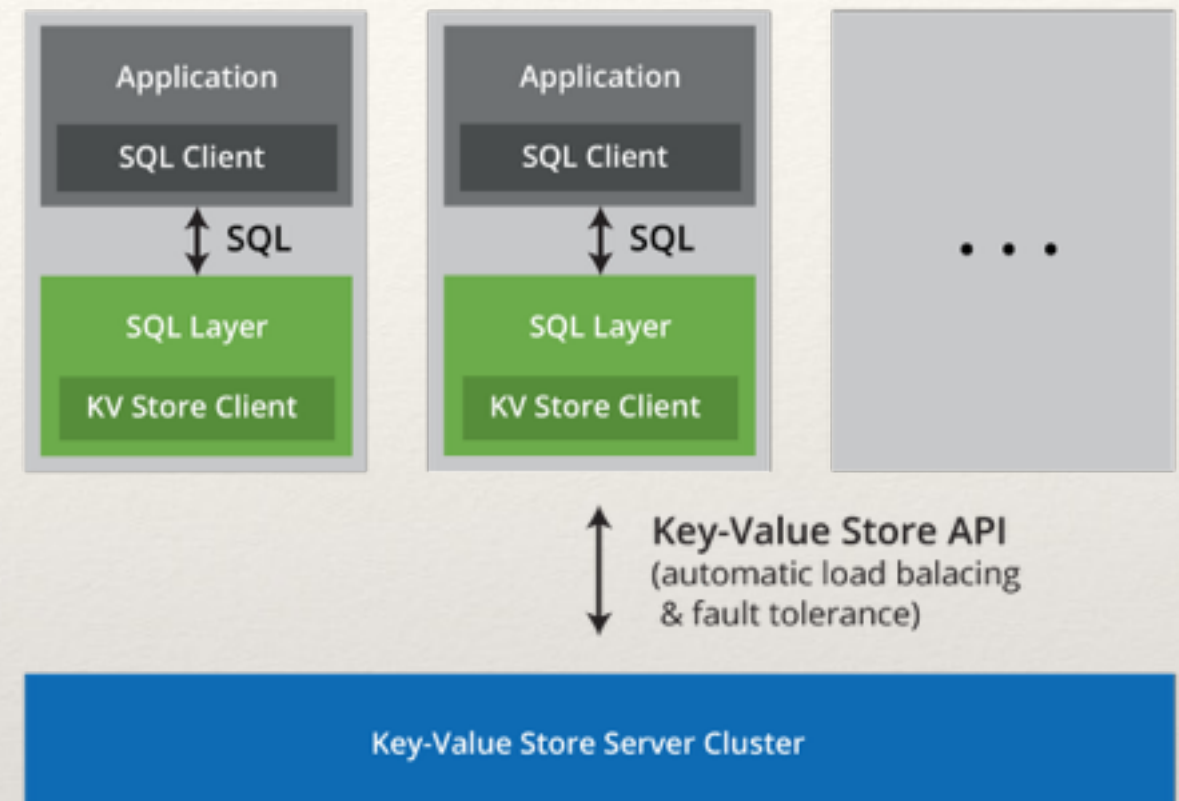


Throughput (per core)



FoundationDB:SQL Layer

- ❖ SQL - layer on top of KV -> transactional, scalable, HA
- ❖ SQL Layer is stateless -> scalable, fault tolerant
- ❖ Hierarchical schema
- ❖ SQL and JSON interfaces
- ❖ Powerful indexing (multi-table, geospatial, ...)

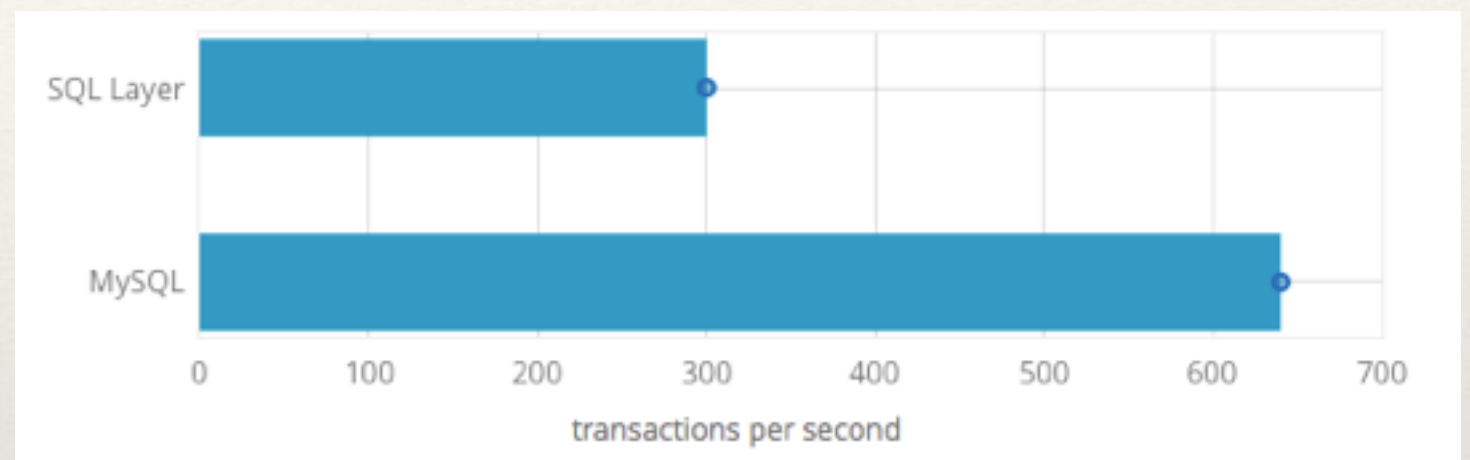


FoundationDB: SQL Performance

Sysbench: read/write, ~80GB, 300M rows

One node test

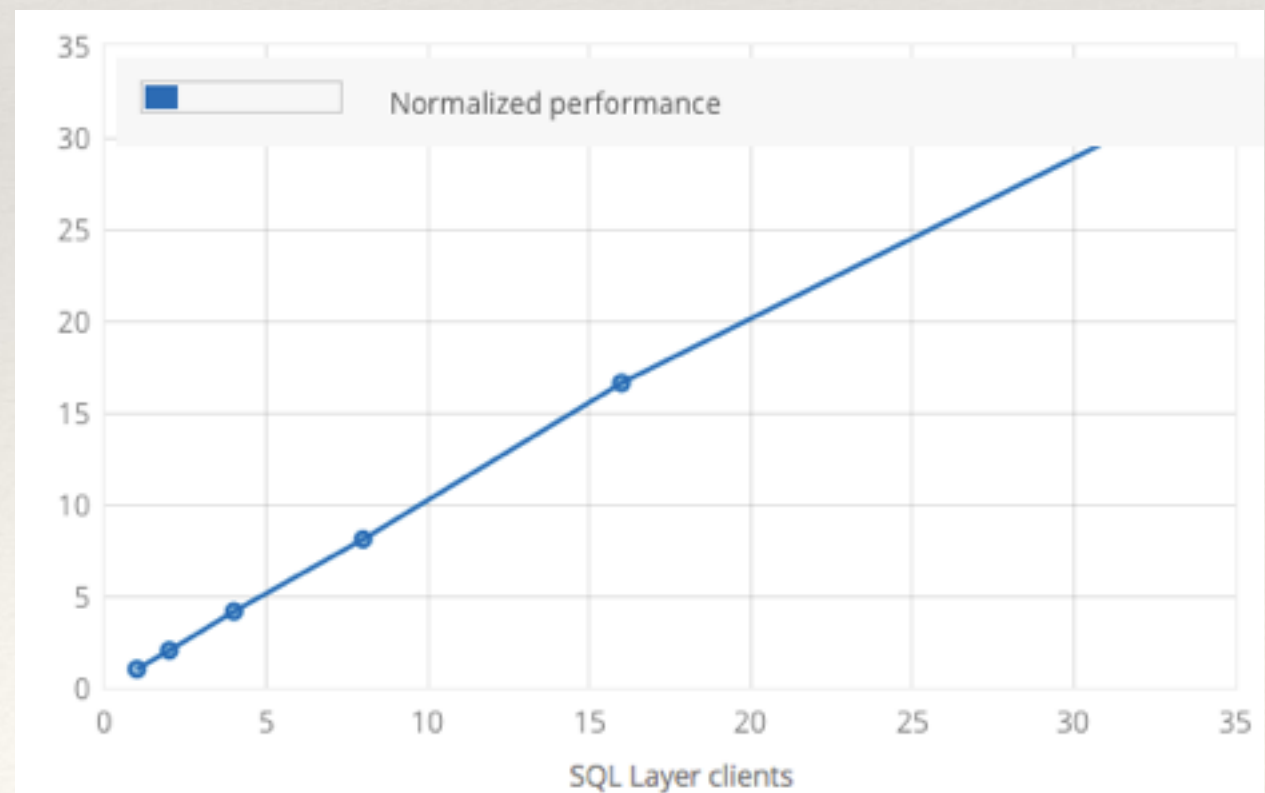
4 core, 16GB RAM, 200GB SATA SSD



Multi nodes test

KV: 8 nodes with 1-process; 3-replication

SQL: up to 32 nodes with
8-thread sysbench process



MemSQL

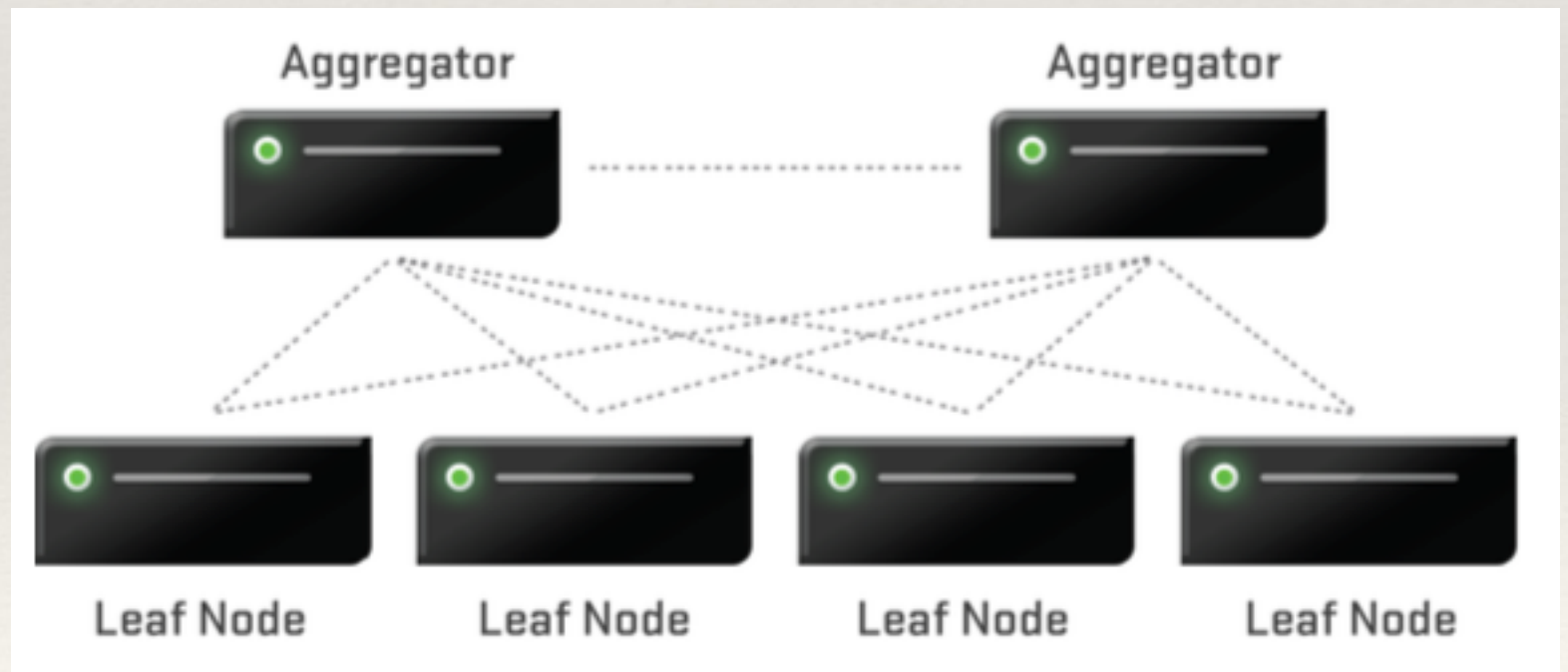


- ❖ In-Memory Storage for OLTP
- ❖ Column-oriented Storage for OLAP
- ❖ Compiled Query Execution Plans (+cache)
- ❖ Local ACID transactions (no global txs for distributed)
- ❖ Lock-free, MVCC
- ❖ Fault tolerance, automatic replication, redundancy (=2 by default)
- ❖ [Almost] no penalty for replica creation

MemSQL



- ❖ Two-tiered shared-nothing architecture
 - Aggregators for query routing
 - Leaves for storage and processing
- ❖ Integration:
 - SQL
 - MySQL protocol
 - JSON API



MemSQL: CAP & ACID

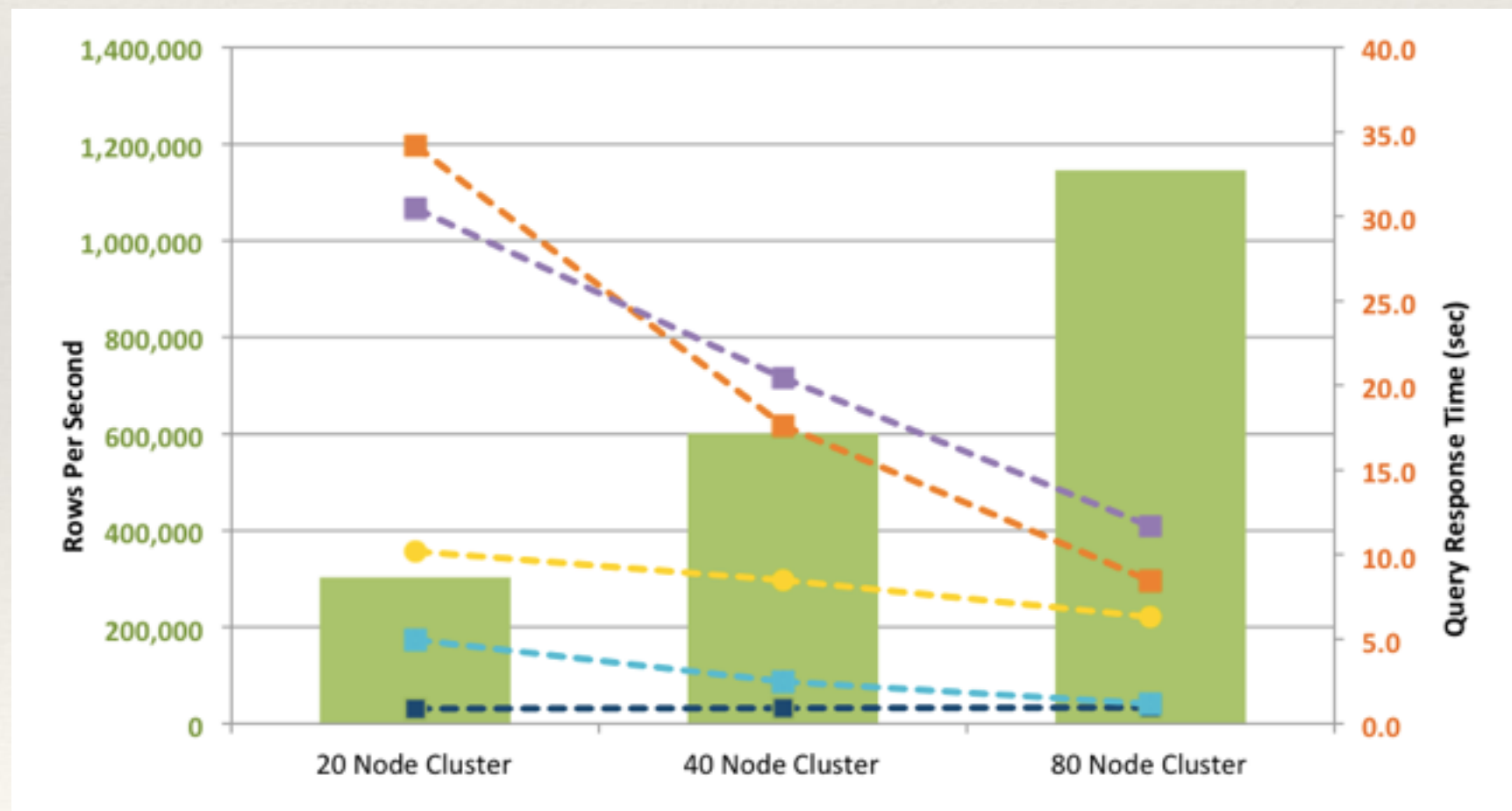


- ❖ `CP` system. Need majority of nodes (or half with master) to work
- ❖ Only 'Read Committed' isolation level ('phantom reads', 'non-repeatable reads' are possible)
- ❖ Manual Master Aggregator management

MemSQL: Performance



- ❖ Adapted TPC-H
- ❖ OLAP Reads & OLTP writes simultaneously
- ❖ AWS EC2 VPC



Overview

	Max Isolation	Scalable	Open Source	Free to try	Language
PostgreSQL	S	Postgres-XL?	Yes	Yes	C
NuoDB	CR	Yes	No	<5 domains	C++
VoltDB	S	Yes	Yes	Yes (wo HA)	Java/C++
ScaleDB	RC?	Yes?	No	?	?
ClustrixDB	RR	Yes	No	Trial (via email req)	C ?
FoundationDB	S	Yes	Partly	<6 processes	Flow(C++)
MemSQL	RC	Yes	No	?	C++

S: Serializable, RR: Read Committed, RC: Read Committed, CR: Consistent Read

Conclusions

- ❖ NewSQL is an established trend with a number of options
- ❖ Hard to pick one because they're not on a common scale
- ❖ No silver bullet
- ❖ Growing data volume requires ever more efficient ways to store and process it

Questions?

Links: General concepts

- ❖ CAP explanation from Brewer, 12 years later
- ❖ Scalable performance, simple explanation
- ❖ What is NewSQL
- ❖ Overview about NoSQL databases
- ❖ Performance loss in OLTP systems
- ❖ Memory price trends
- ❖ (wiki) Shared Nothing Architecture
- ❖ (wiki) Column oriented DBMS
- ❖ How NewSQL handles big data
- ❖ What is YCSB benchmark
- ❖ What is TPC benchmark
- ❖ Transactional isolation levels

Links: NuoDB

- ❖ <http://www.infoq.com/articles/nuodb-architecture-1/>
- ❖ <http://www.infoq.com/articles/nuodb-architecture-2/>
- ❖ <http://stackoverflow.com/questions/14552091/nuodb-and-hdfs-as-storage>
- ❖ http://go.nuodb.com/rs/nuodb/images/NuoDB_Benchmark_Report.pdf
- ❖ NuoDB white paper (google has you :)
- ❖ <https://aphyr.com/posts/292-call-me-maybe-nuodb>
- ❖ <http://dev.nuodb.com/techblog/failure-detection-and-network-partition-management-nuodb>

Links: VoltDB

- ❖ White paper, Technical overview (google has you)
- ❖ <https://github.com/VoltDB/voltdb-client-erlang/blob/master/doc/BENCHMARK1.md>
- ❖ <http://www.mysqlperformanceblog.com/2011/02/28/is-voltdb-really-as-scalable-as-they-claim/>
- ❖ <https://voltdb.com/blog/voltdb-3-x-performance-characteristics/>
- ❖ <http://docs.voltdb.com/UsingVoltDB/KsafeNetPart.php>
- ❖ <https://news.ycombinator.com/item?id=6639127>

Links: ScaleDB

- ❖ <http://scaledb.com/pdfs/TechnicalOverview.pdf>
- ❖ http://www.scaledb.com/pdfs/scaledb_multitenant.pdf
- ❖ http://www.percona.com/live/mysql-conference-2013/sites/default/files/slides/DB_Vistualization_for_PublicPrivate_Clouds.pdf

Links: Clustrix

- ❖ http://www.clustrix.com/wp-content/uploads/2013/10/Clustrix_A-New-Approach_WhitePaper.pdf
- ❖ http://www.clustrix.com/wp-content/uploads/2013/10/Clustrix_Driving-the-New-Wave_WP.pdf
- ❖ http://www.clustrix.com/wp-content/uploads/2013/10/Clustrix_AWS_WP.pdf
- ❖ http://www.clustrix.com/wp-content/uploads/2013/10/Clustrix_TPCC_Percona.pdf
- ❖ <http://sergei.clustrix.com/2011/01/mongodb-vs-clustrix-comparison-part-1.html>
- ❖ <http://docs.clustrix.com/display/CLXDOC/Consistency%2C+Fault+Tolerance%2C+and+Availability>

Links: FoundationDB

- ❖ <https://foundationdb.com/key-value-store/white-papers>
- ❖ <http://blog.foundationdb.com/call-me-maybe-foundationdb-vs-jepsen>
- ❖ <https://foundationdb.com/acid-claims>
- ❖ <https://foundationdb.com/key-value-store/performance>
- ❖ <https://foundationdb.com/layers/sql/documentation/Concepts>
- ❖ <https://foundationdb.com/layers/sql/documentation/SQL/indexes.html>
- ❖ <https://foundationdb.com/layers/sql/performance>
- ❖ <https://foundationdb.com/key-value-store/features>
- ❖ <https://foundationdb.com/key-value-store/documentation/configuration.html>
- ❖ <https://foundationdb.com/key-value-store/documentation/beta1/developer-guide.html>
- ❖ <https://foundationdb.com/layers/sql/documentation/Concepts/known.limitations.html>

Links: MemSQL

- ❖ MemSQL Whitepaper "The Modern Database Landscape"
- ❖ MemSQL Whitepaper "ESG Lab Benchmark of MemSQL's Performance"
- ❖ MemSQL Whitepaper "Technical overview"
- ❖ http://developers.memsql.com/docs/latest/concepts/dev_concepts.html
- ❖ http://developers.memsql.com/docs/2.6/admin/high_availability.html